

Conditional Disclosure of Secrets via Non-Linear Reconstruction

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Conditional Disclosure of Secrets

[Gertner-Ishai-Kushilevitz-Malkin'00]

A



$f : [N] \rightarrow \{0, 1\}$

B



$i \in [N]$

C



f, i

Conditional Disclosure of Secrets

[Gertner-Ishai-Kushilevitz-Malkin'00]

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$f : [N] \rightarrow \{0, 1\}$
bit *s*

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$i \in [N]$
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C

f, i

Conditional Disclosure of Secrets

[Gertner-Ishai-Kushilevitz-Malkin'00]



$f : [N] \rightarrow \{0, 1\}$
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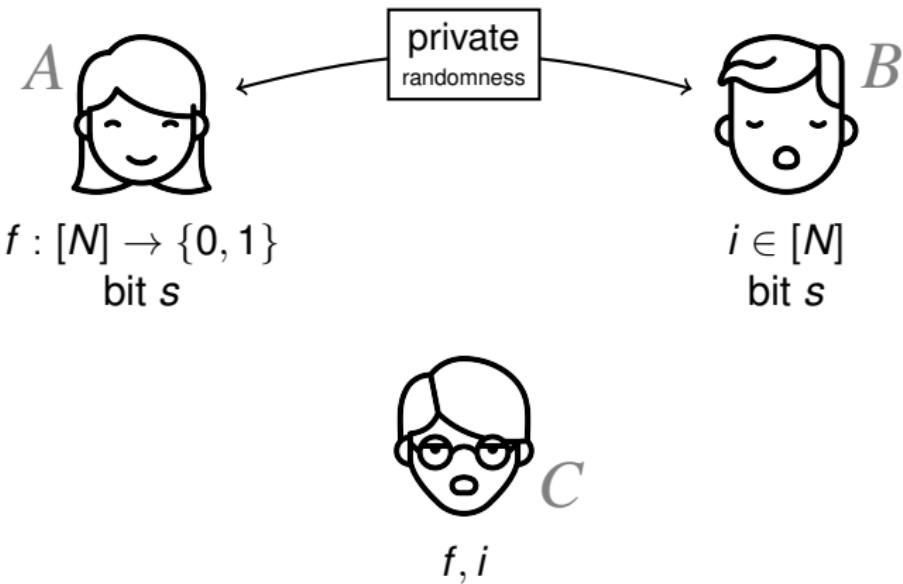


f, i

Charlie gets s iff $f(i) = 1$

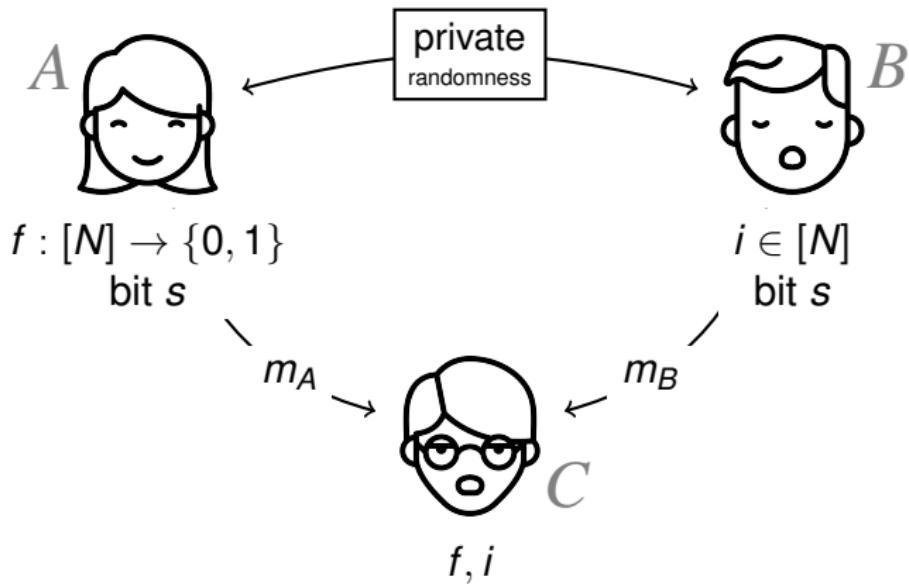
Conditional Disclosure of Secrets

[Gertner-Ishai-Kushilevitz-Malkin'00]



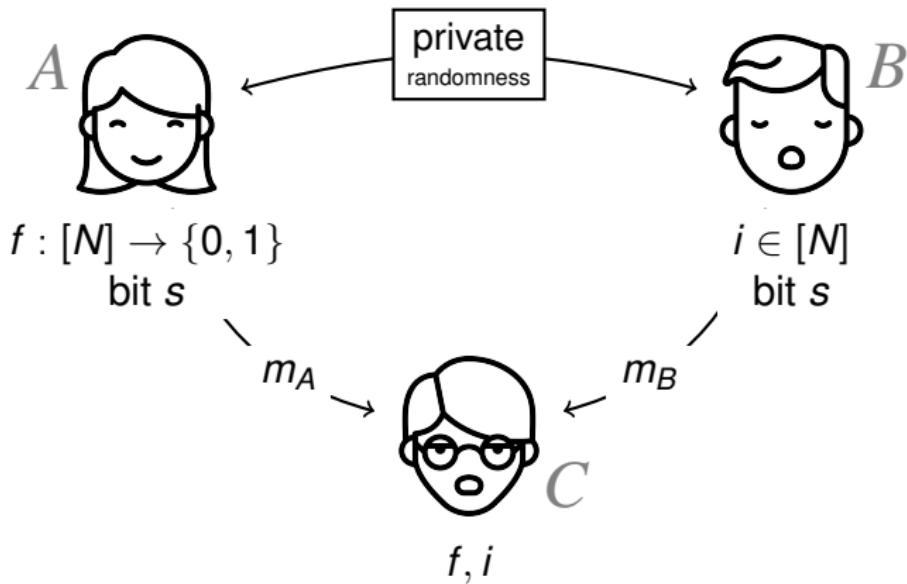
Conditional Disclosure of Secrets

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Conditional Disclosure of Secrets

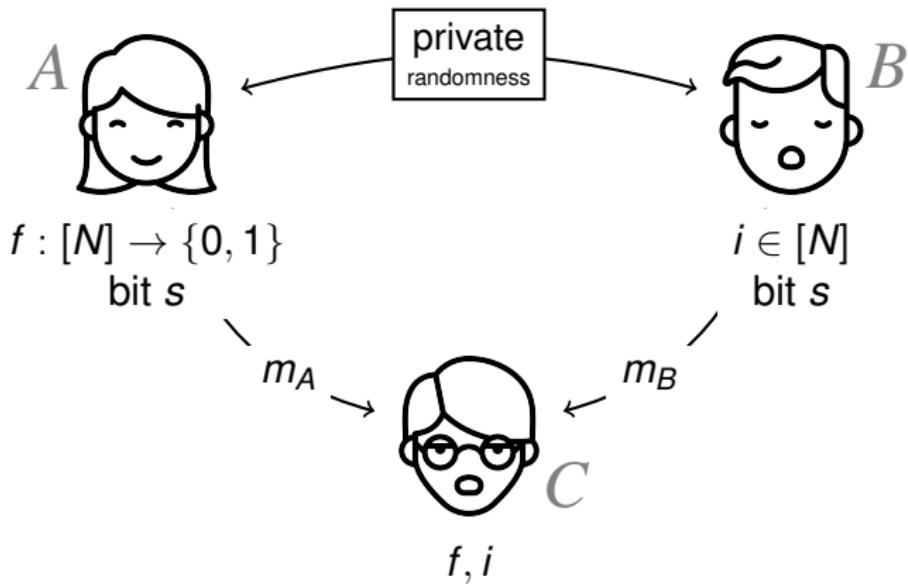
[Gertner-Ishai-Kushilevitz-Malkin'00]



- ▶ Correctness: When $f(i) = 1$, Charlie gets s .

Conditional Disclosure of Secrets

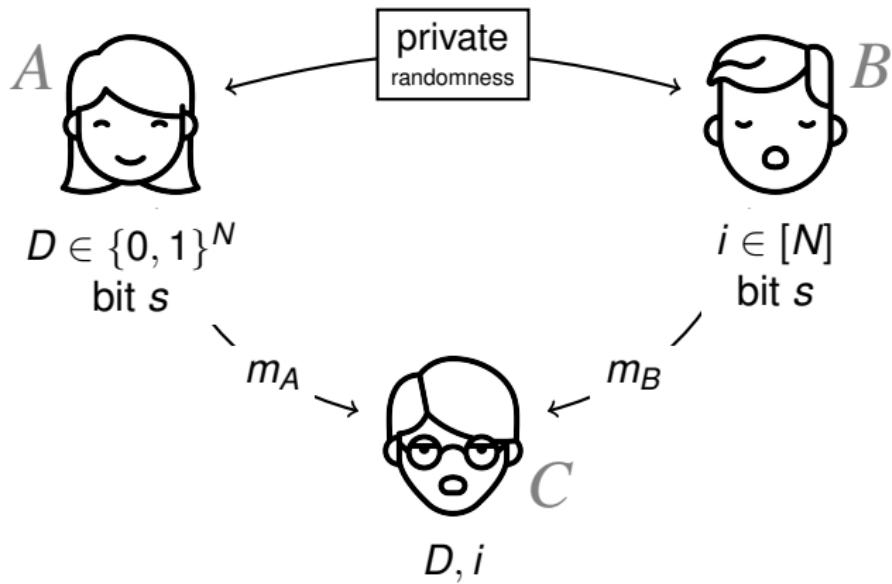
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- ▶ Correctness: When $f(i) = 1$, Charlie gets s .
- ▶ IT Privacy: When $f(i) = 0$, (m_A, m_B) can be *perfectly simulated*.

Conditional Disclosure of Secrets

[Gertner-Ishai-Kushilevitz-Malkin'00]



- ▶ Correctness: When $D_i = 1$, Charlie gets s .
- ▶ IT Privacy: When $D_i = 0$, (m_A, m_B) can be *perfectly simulated*.

Conditional Disclosure of Secrets

- ▶ Symmetric PIR [GIKM'00]
 - ▶ How to handle malicious clients in SPIR
- ▶ Secret Sharing [SS'97,BIKK'14]
 - ▶ For certain graph-based access structures
- ▶ Attribute-Based Encryption [Att'14,Wee'14,CGW'15]
 - ▶ CDS = 1-key, 1-ciphertext, private-key ABE

Simple CDS Protocol I



$D \in \{0, 1\}^N$
bit s

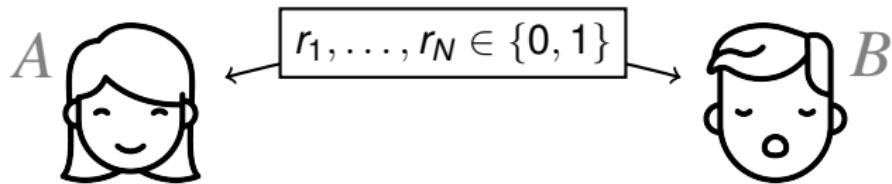


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Simple CDS Protocol I

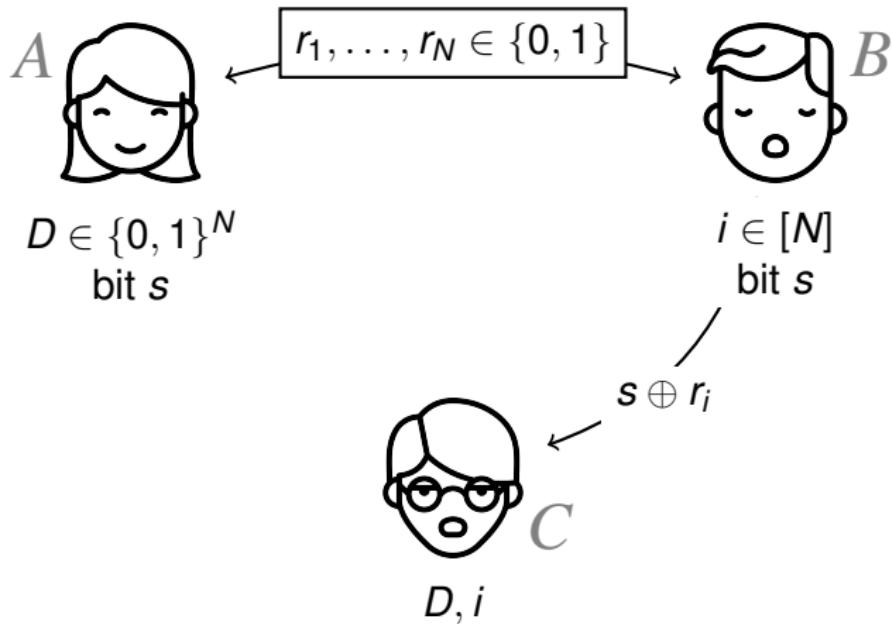


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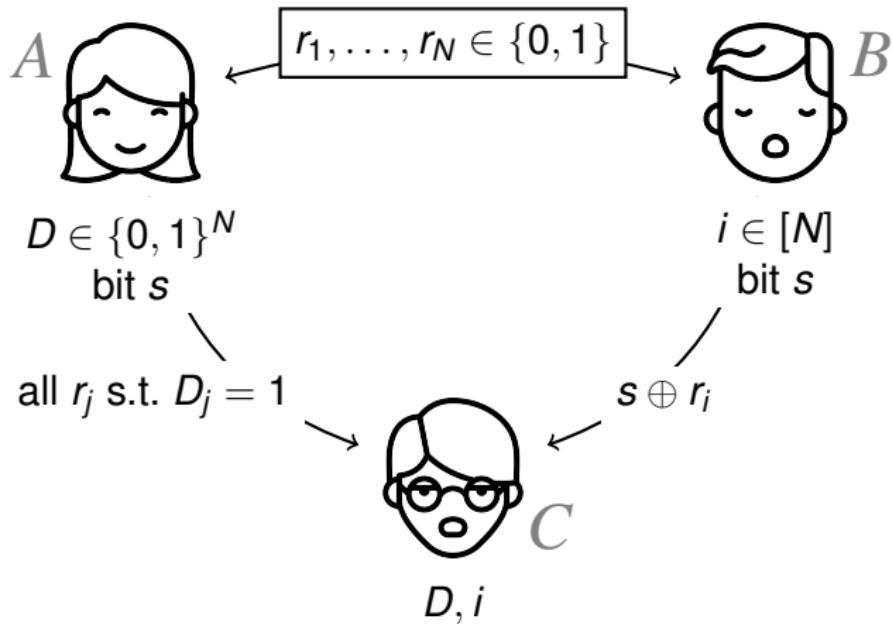
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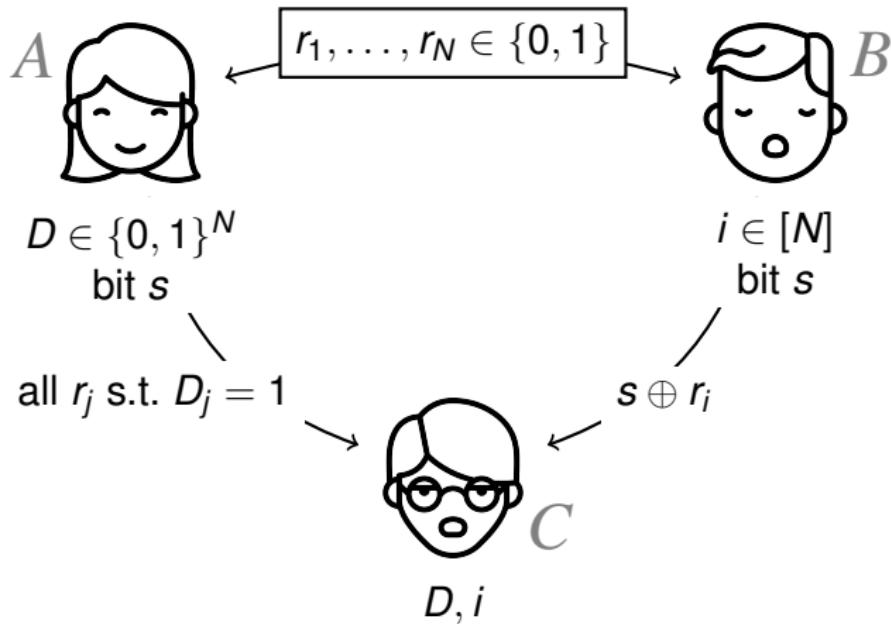
Simple CDS Protocol I



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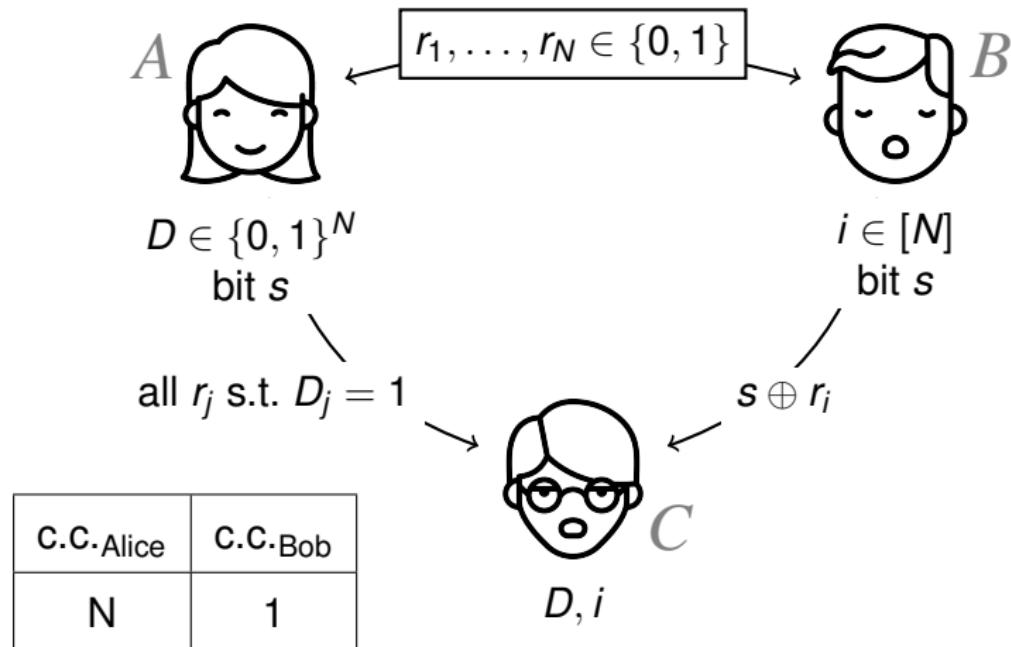


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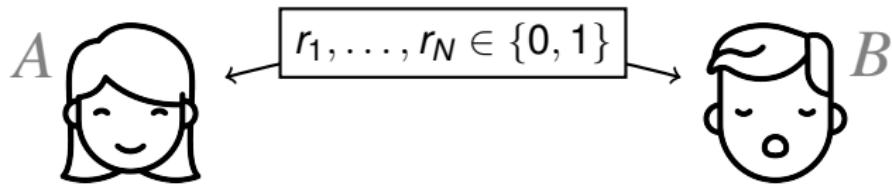
- ▶ Correctness: When $D_i = 1$, Alice sends r_i
- ▶ IT Privacy: When $D_i = 0$, Alice does not send r_i

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Simple CDS Protocol II

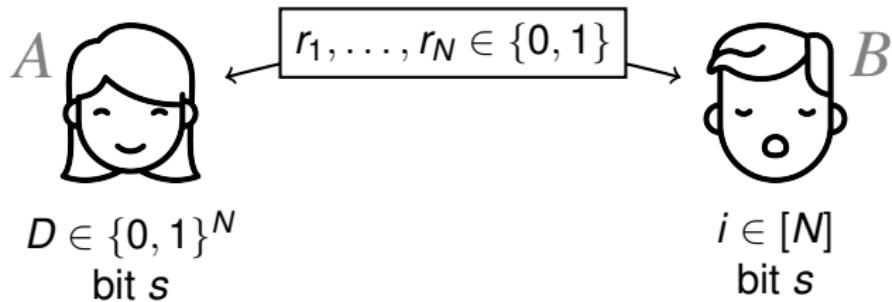


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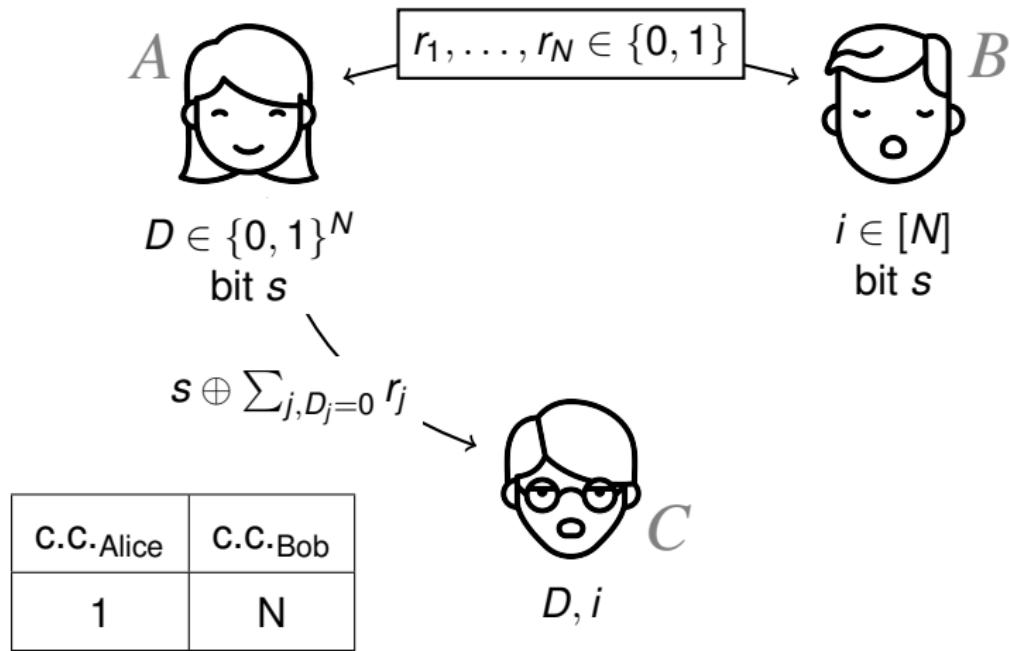
Simple CDS Protocol II



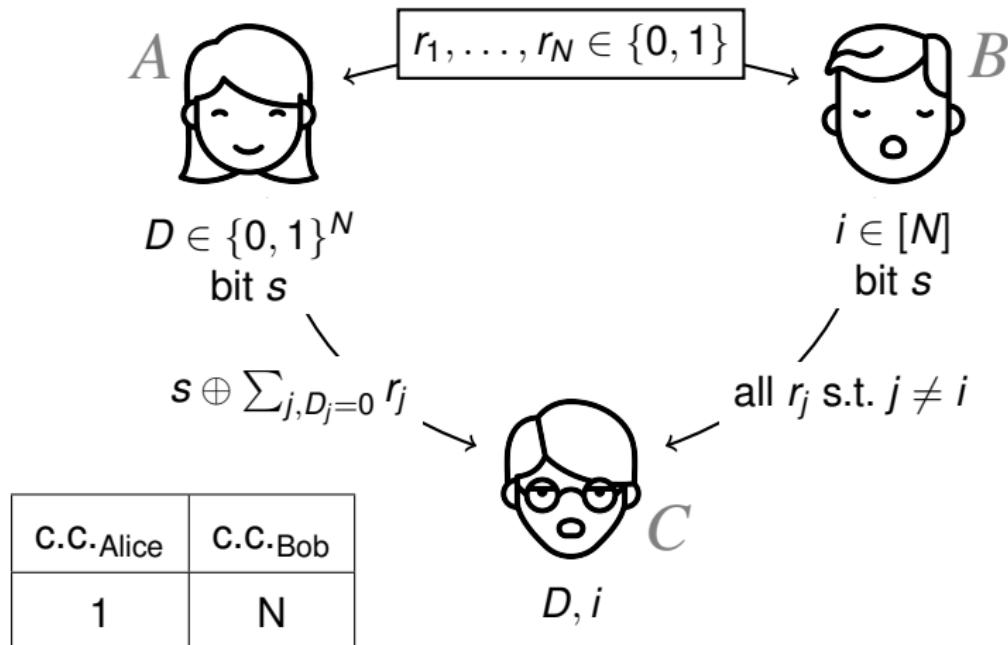
C.C. Alice	C.C. Bob
1	N



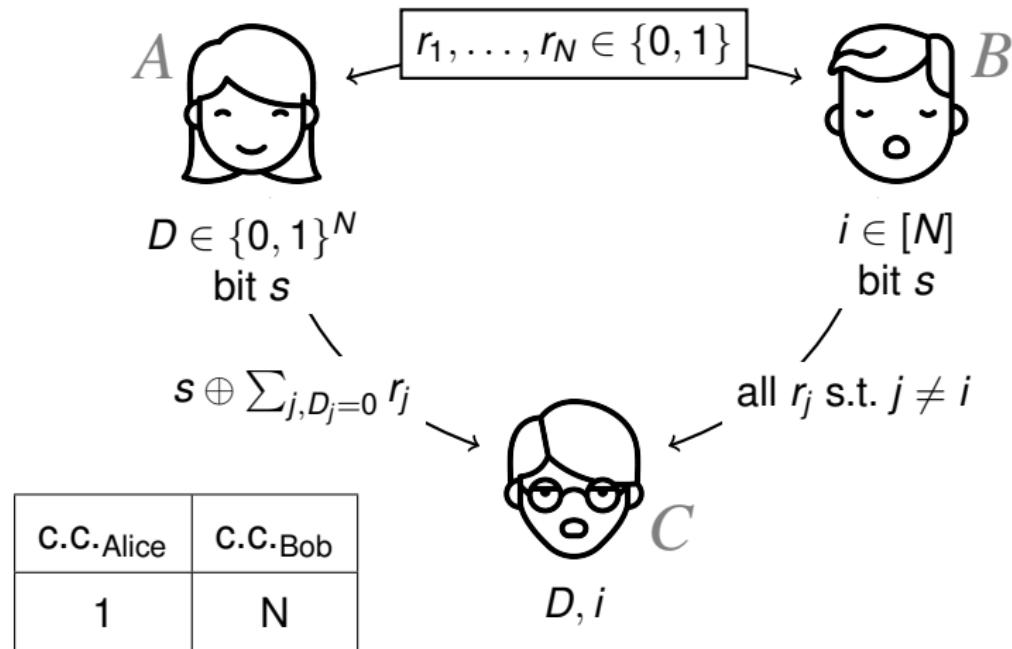
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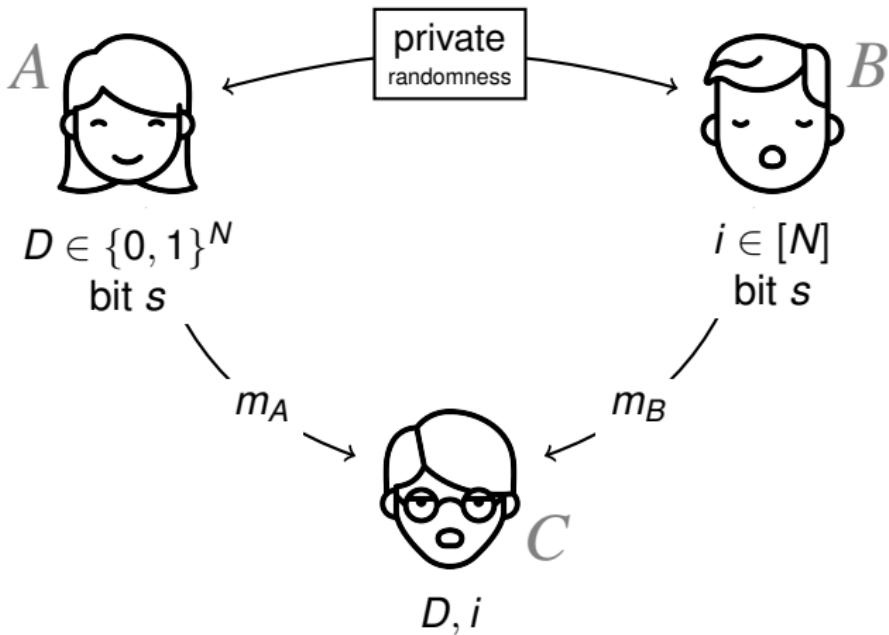


Simple CDS Protocol II



- ▶ Correctness: When $D_i = 1$, Bob sends all r_j s.t. $D_j = 0$
- ▶ IT Privacy: When $D_i = 0$, Bob does not send r_i

CDS with Linear Reconstruction [GKW'15]



Charlie's reconstruction function $C_{i,D}(m_A, m_B) \mapsto s$ is linear.

Previous Works

Communication Complexity	Reconstruction
$(N, 1)$	linear
$(1, N)$	linear

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$\Omega(\sqrt{N})$	[GKW'15]

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$\Omega(\log N)$	[GKW'15]
	general

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Communication Complexity	Reconstruction
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$(1, N)$	linear
----------	--------

$O(\sqrt{N})$	[GKW'15, ...]	linear
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$\Omega(\sqrt{N})$	[GKW'15]	linear
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An exponential gap here
Non-linear reconstruction needed!

$\Omega(\log N)$	[GKW'15]	general
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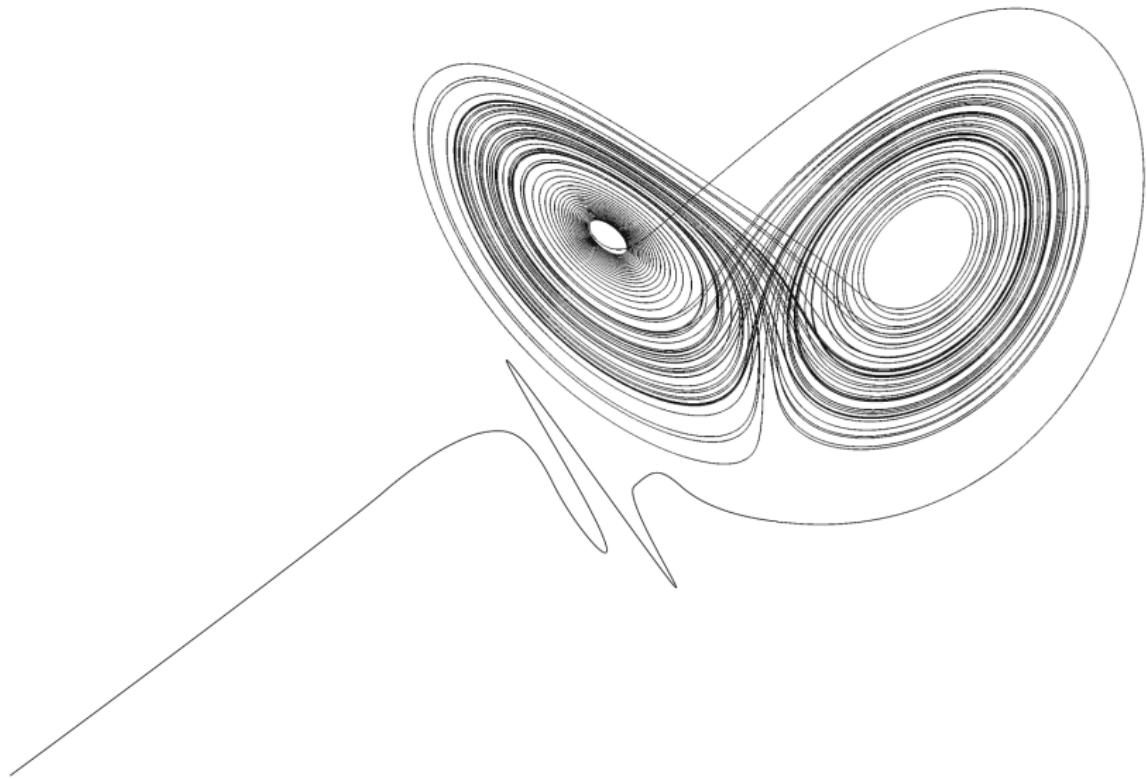
Our Results

Communication Complexity	Reconstruction
$(N, 1)$	linear
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$O(\sqrt{N})$	[GKW'15, ...]
$\Omega(\sqrt{N})$	[GKW'15]
$\Theta(\sqrt[3]{N})$	[This work]
$\Omega(\log N)$	quadratic
$\Omega(\log N)$	[GKW'15]
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CDS: Need Non-linear Techniques!



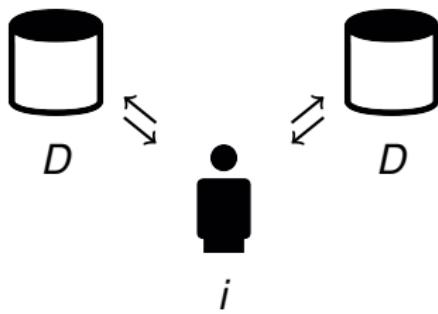
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2-server PIR
[CGKS'95]

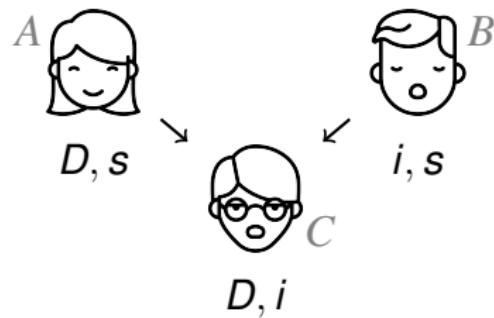
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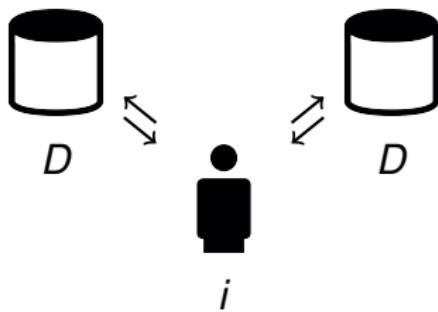


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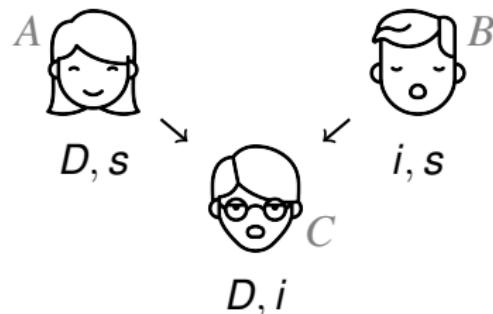
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- ▶ $O(\sqrt{N})$ c.c. [CGKS'95]
linear server

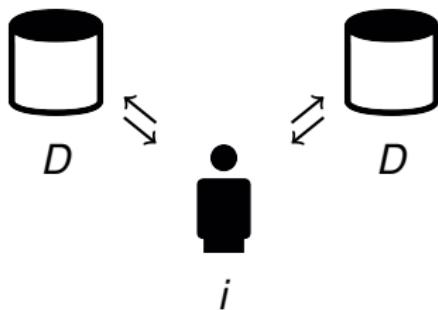
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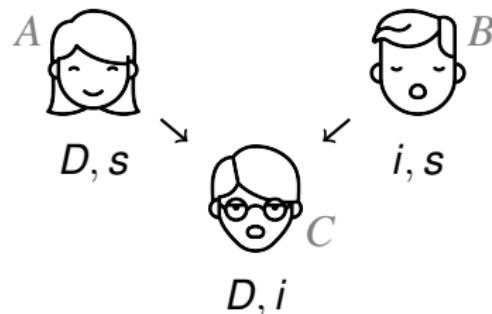
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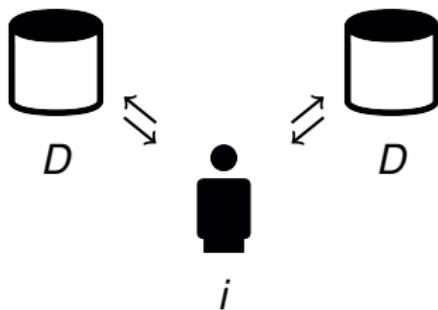


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- ▶ $O(\sqrt[3]{N})$ c.c. [CGKS'95,WY'05]
quadratic server
- ▶ $2^{\tilde{O}(\sqrt{\log N})}$ c.c. [DG'15]
general server

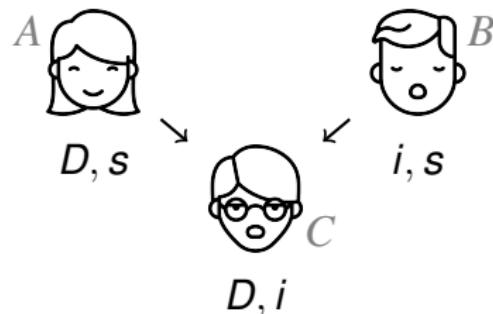
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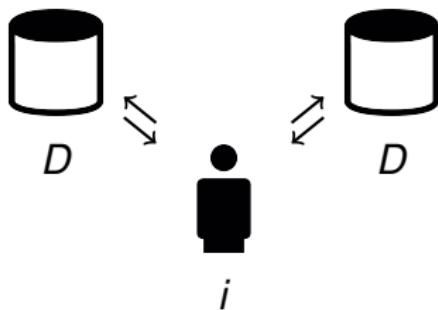
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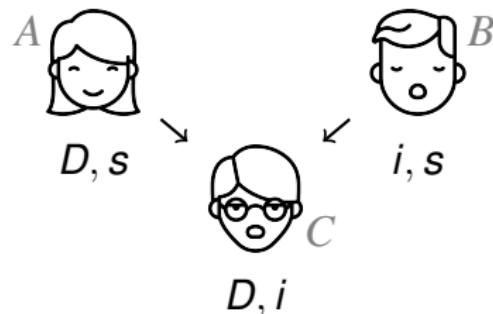
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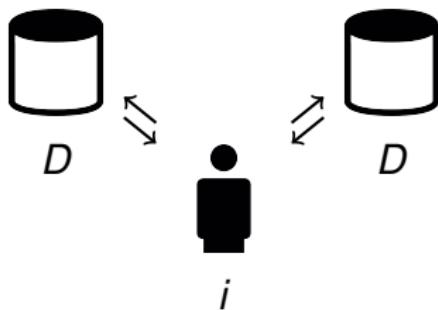
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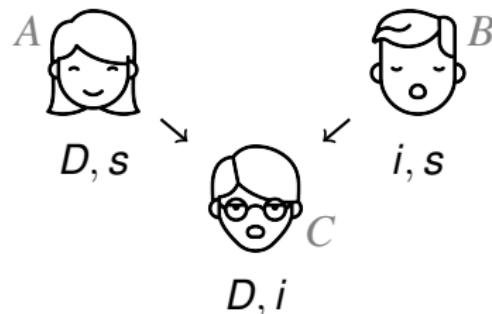
\sqrt{N} c.c. PSM [BIKK'14] from 4-server $\sqrt[4]{N}$ c.c. PIR [CGKS'95]

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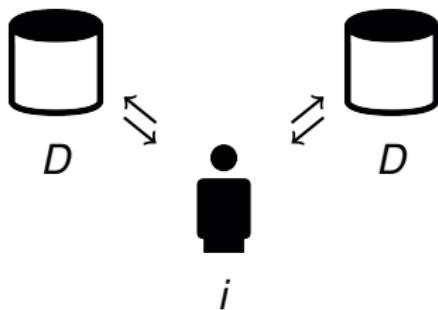
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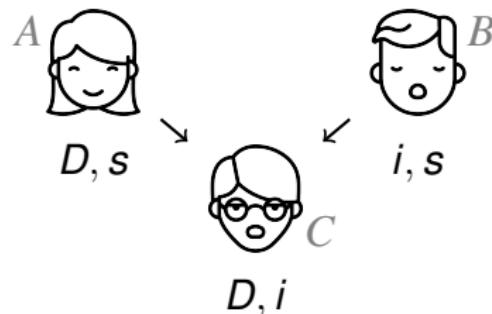
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Private Information Retrieval



$$D \in \{0, 1\}^N$$

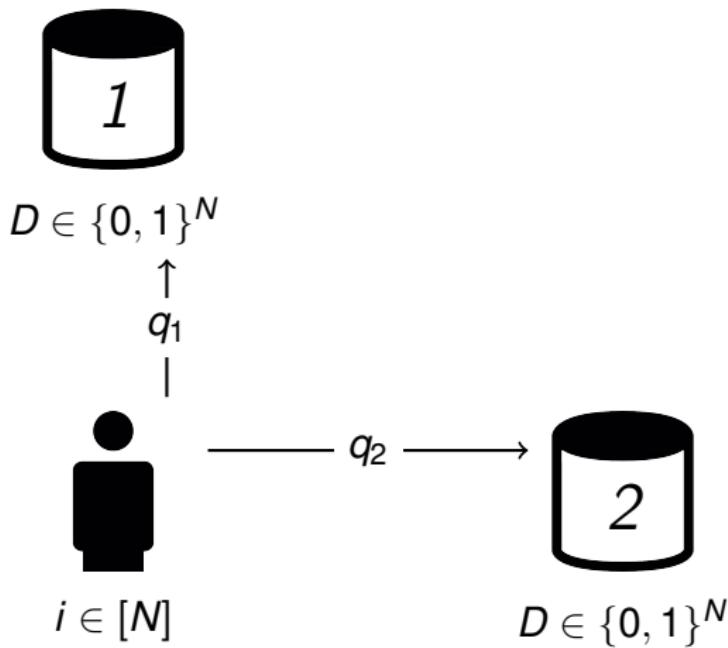


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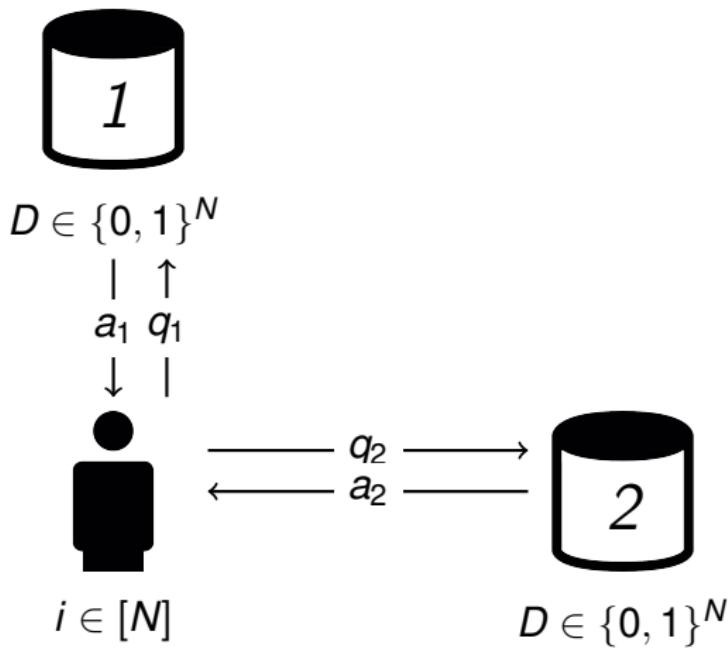


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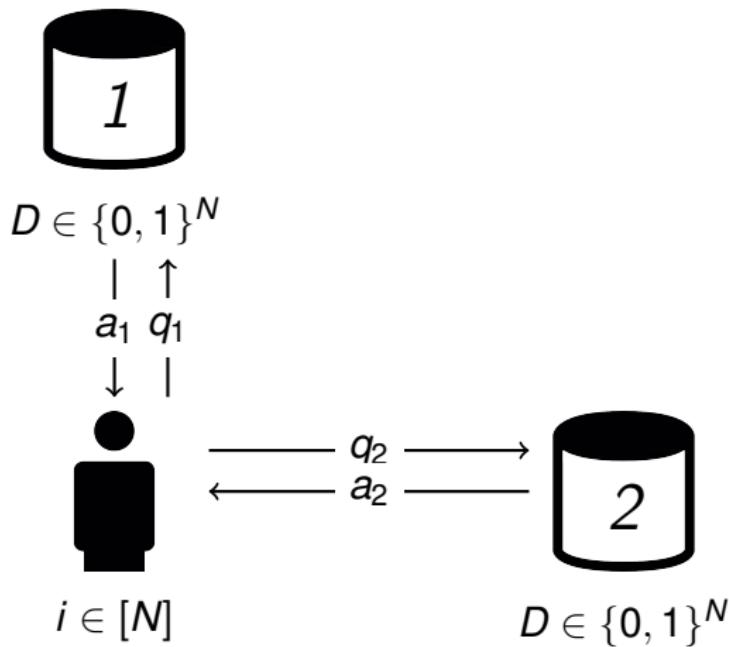
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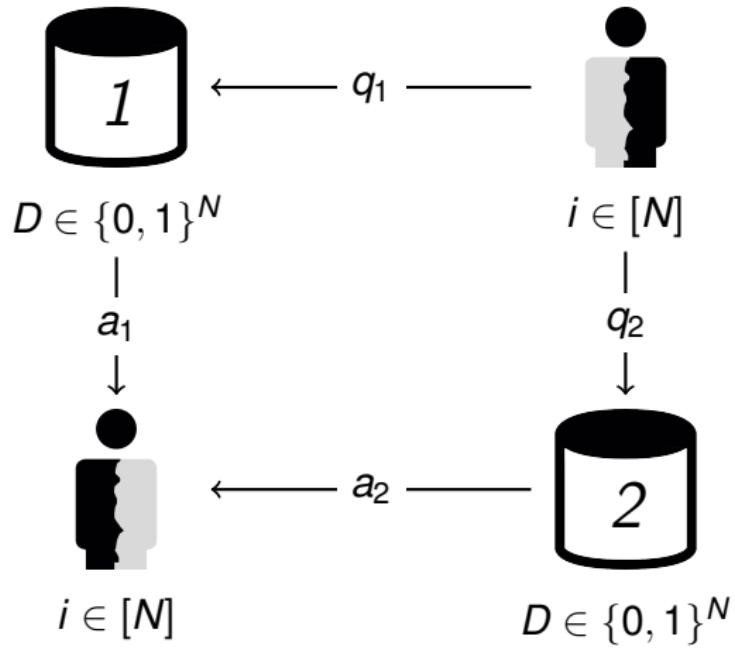


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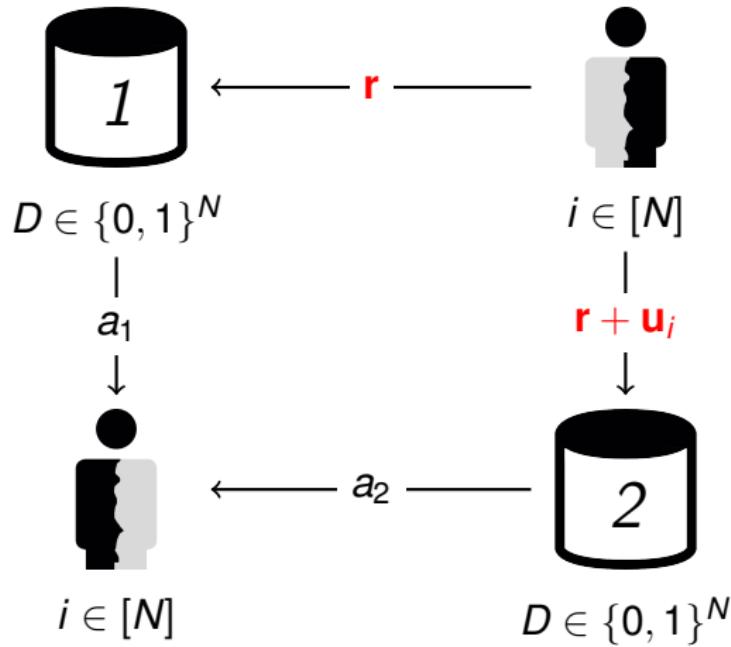


- ▶ Correctness: Client gets D_i
- ▶ IT Privacy: q_1 (resp. q_2) leaks nothing about i

Private Information Retrieval

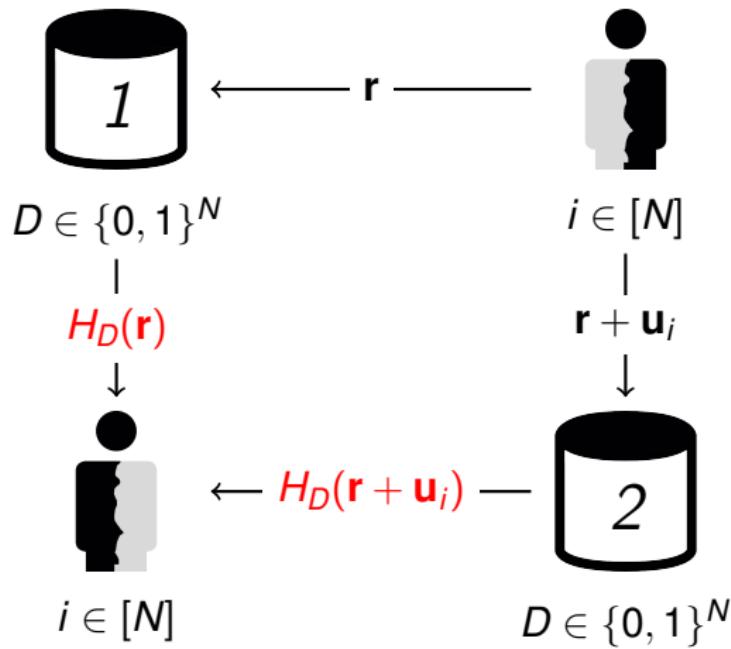


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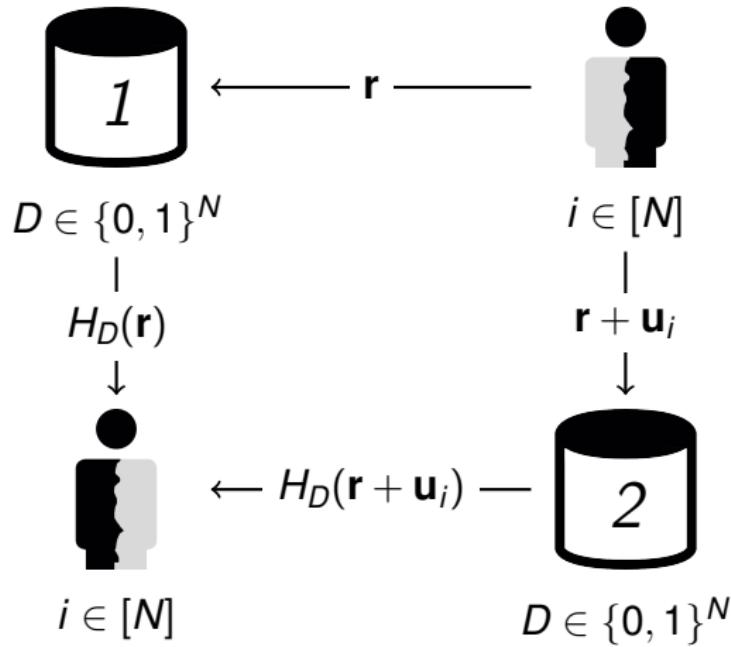
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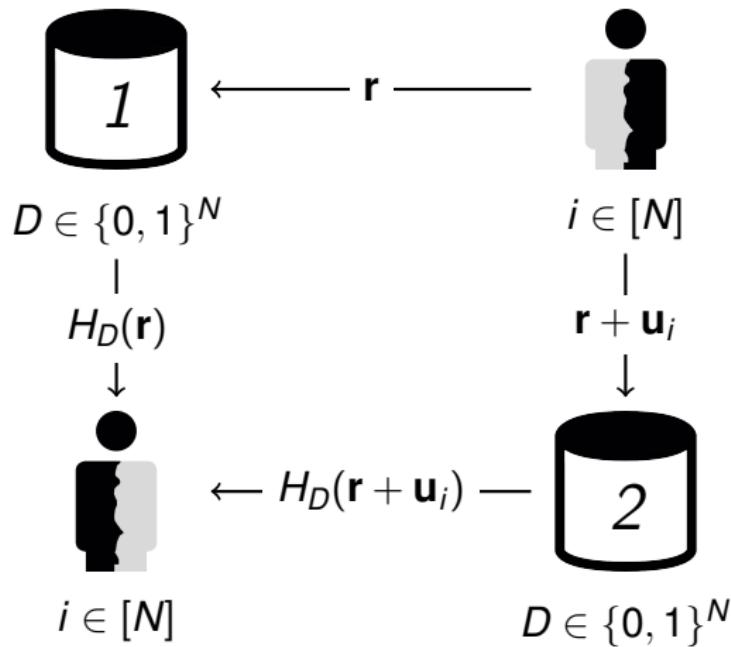
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Prop. 2 “Linear client”: $\langle \mathbf{u}_i, H_D(\mathbf{r} + \mathbf{u}_i) \rangle - \langle \mathbf{u}_i, H_D(\mathbf{r}) \rangle = D_i$

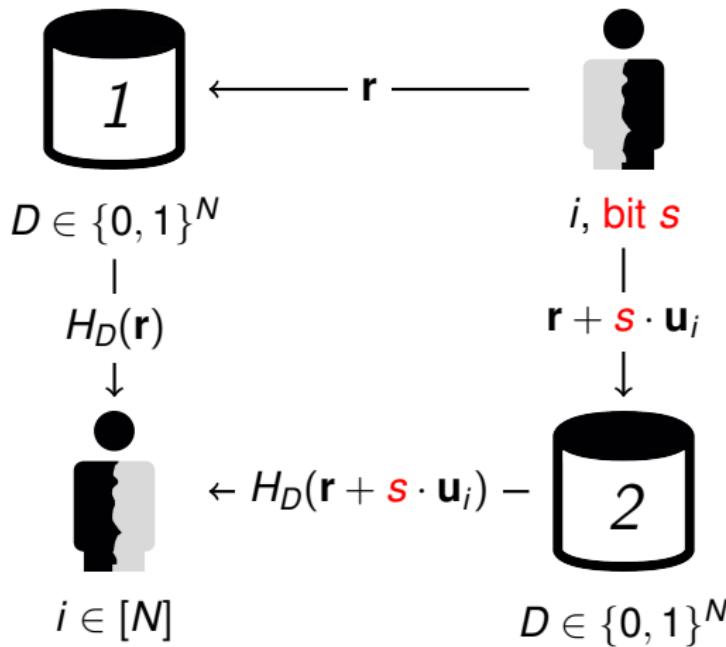
Private Information Retrieval \implies CDS



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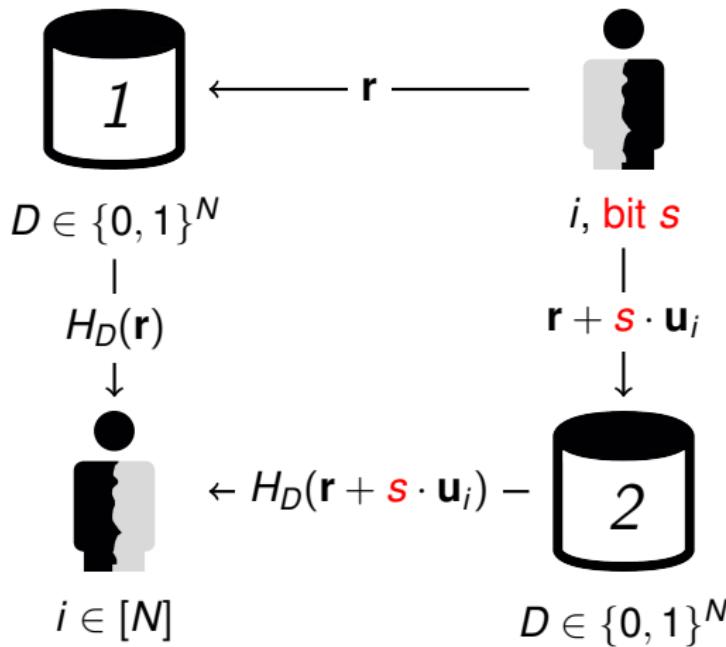
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Prop. 1 Vectors u_1, \dots, u_N . Queries are additive secret sharing of u_i

Prop. 2 “Linear client”: $\langle u_i, H_D(r + s \cdot u_i) \rangle - \langle u_i, H_D(r) \rangle = ?$

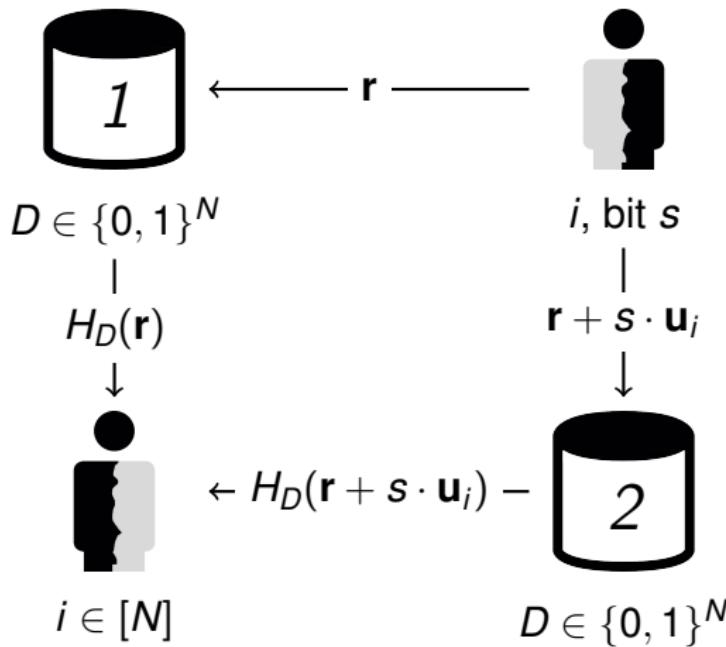
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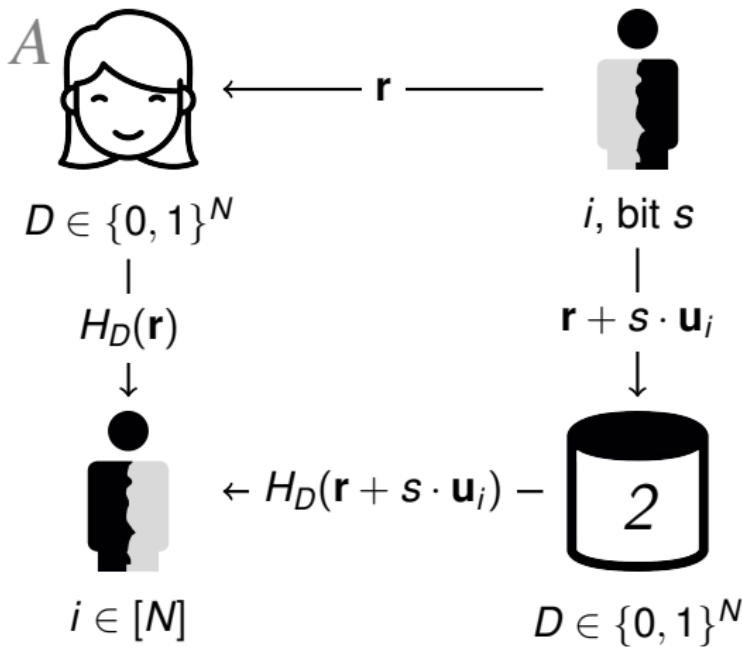
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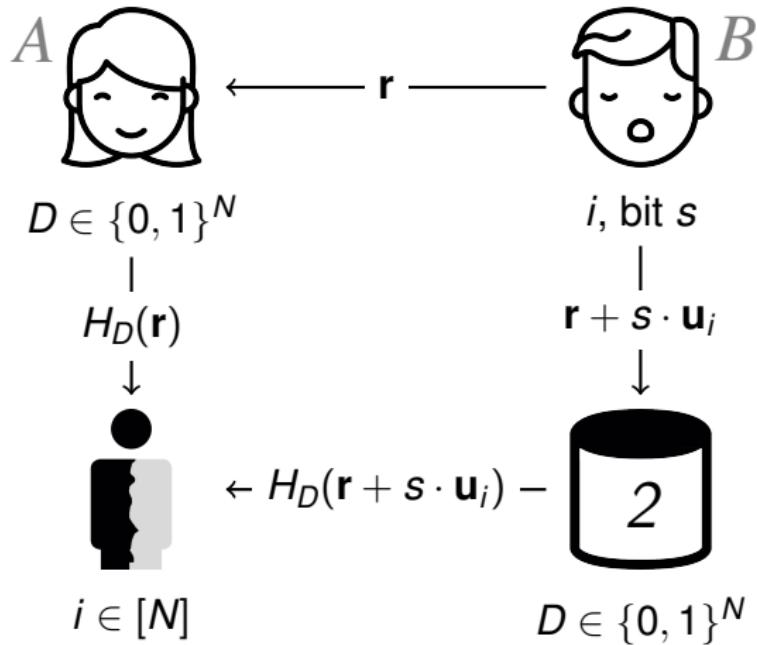
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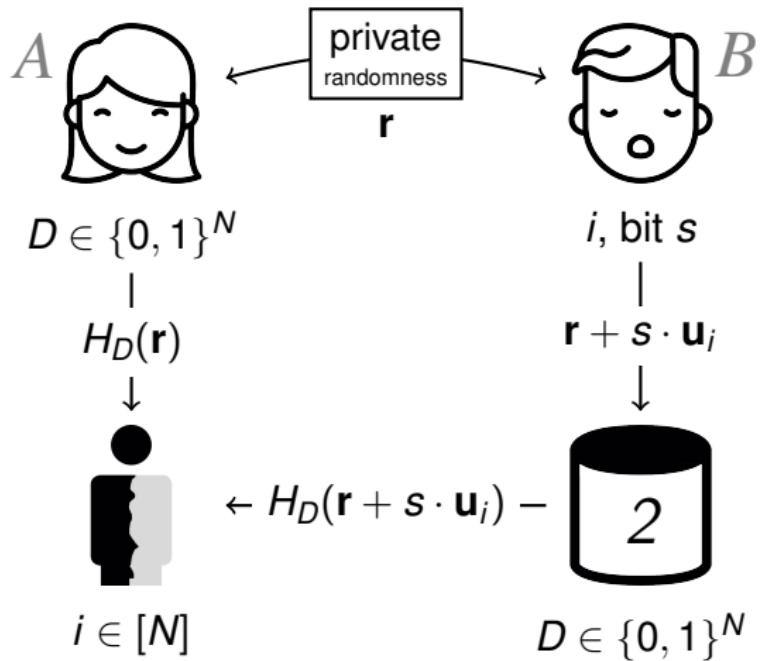
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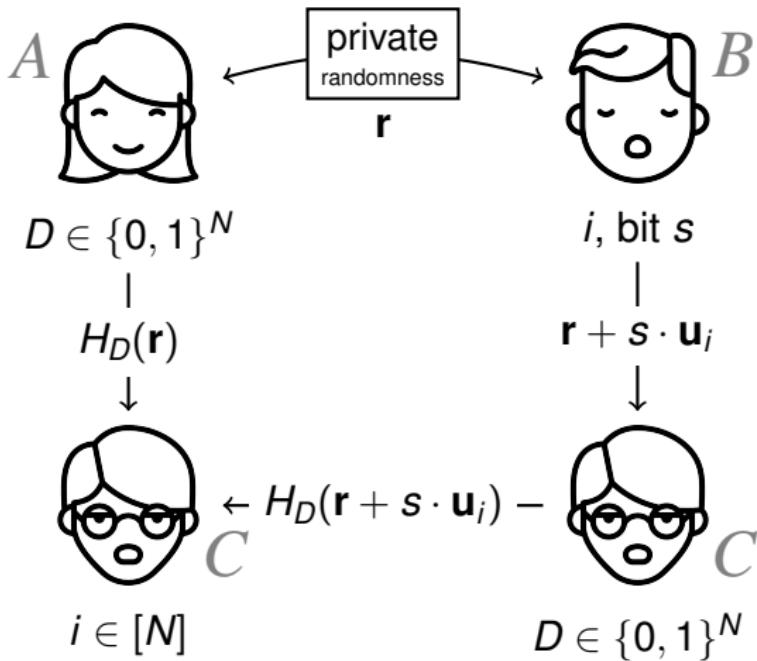
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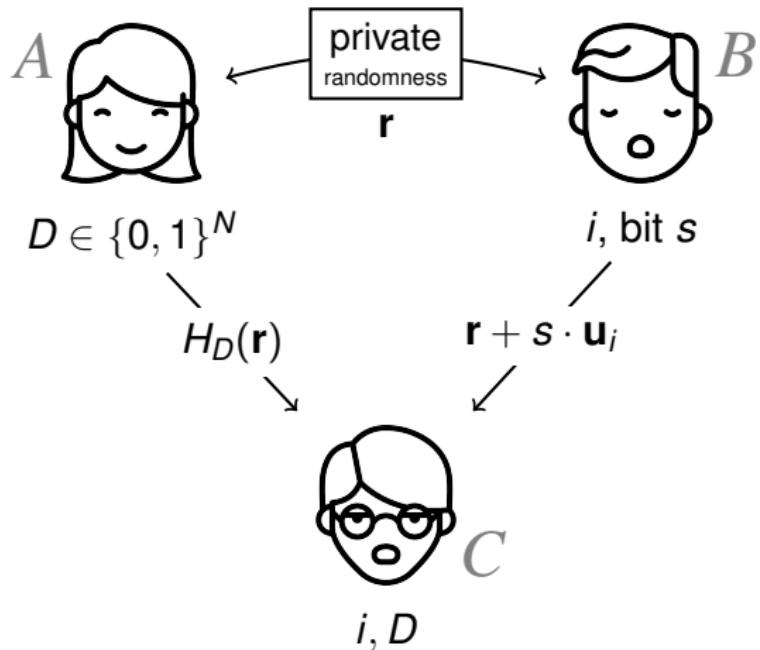
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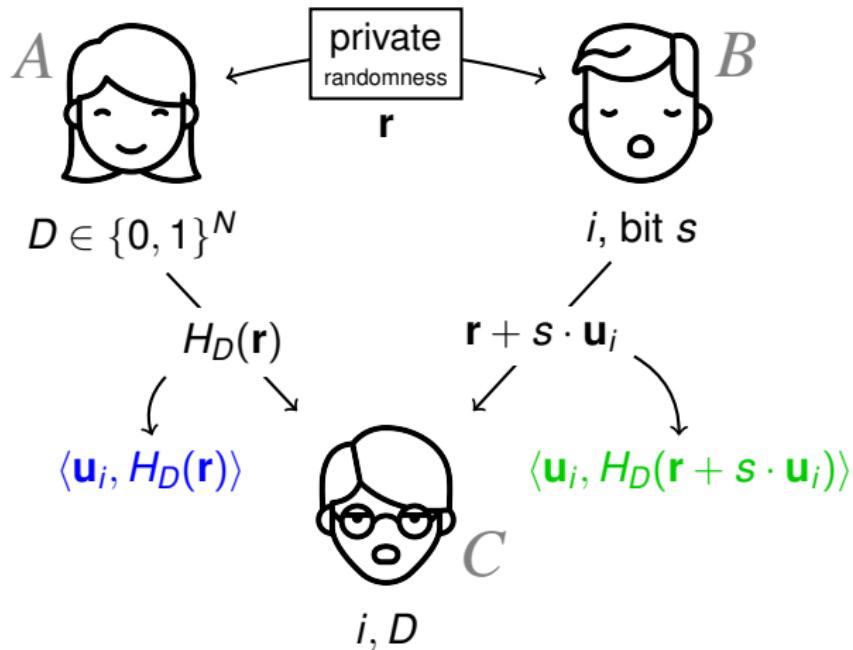
Private Information Retrieval \implies CDS



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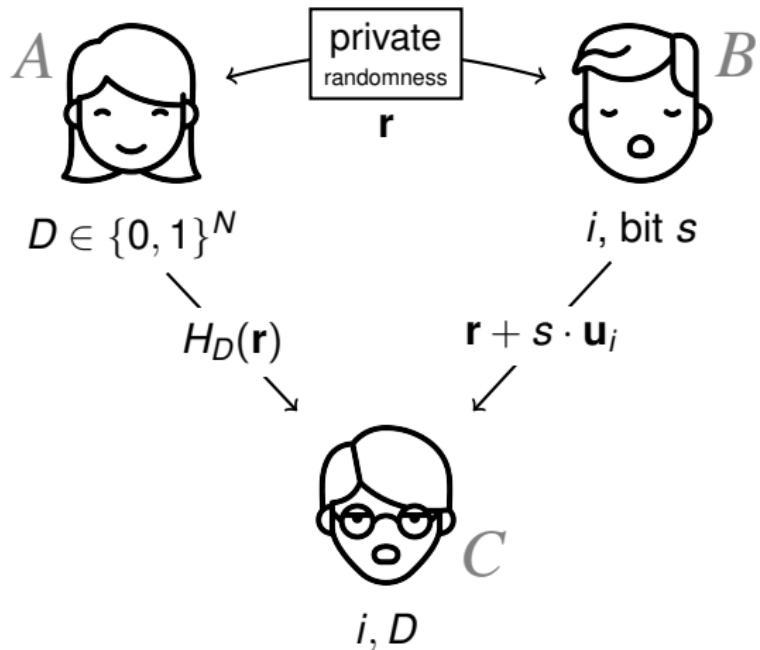
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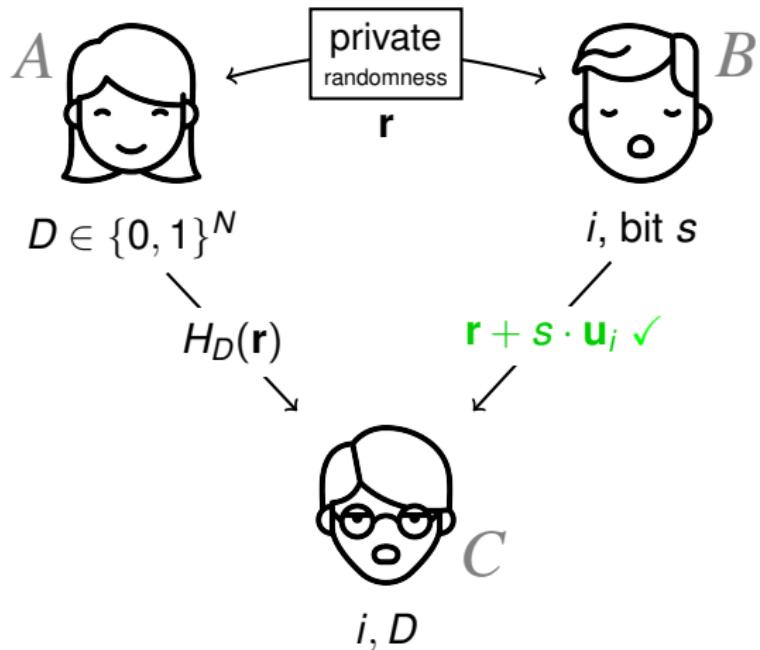
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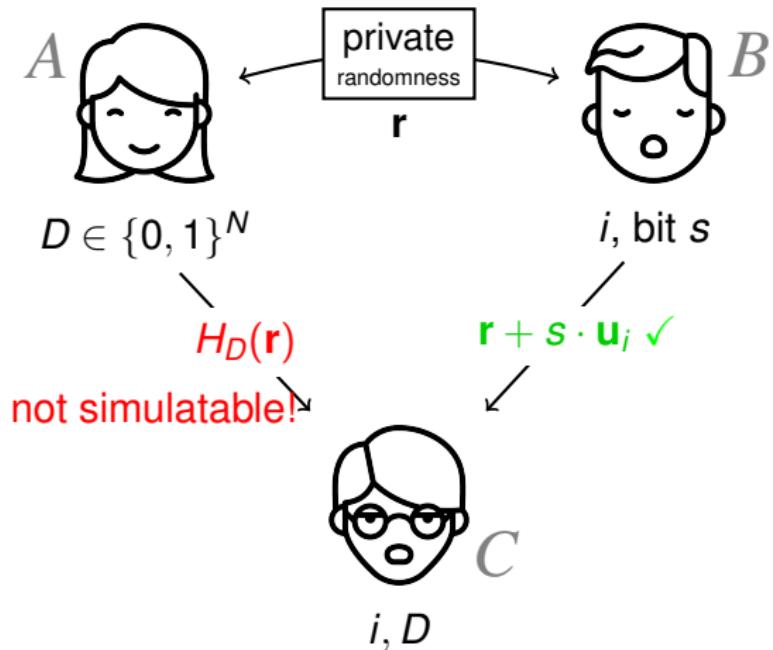
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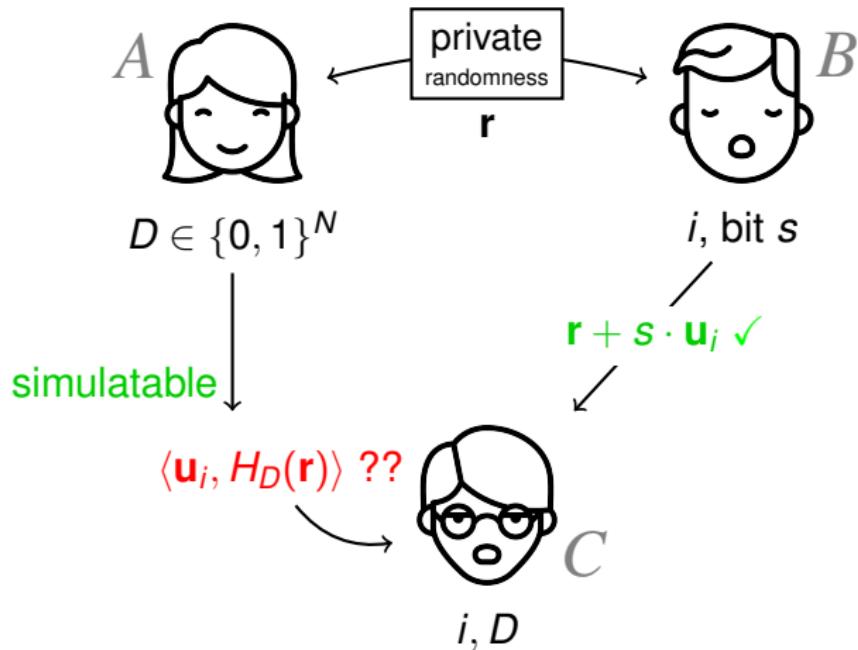
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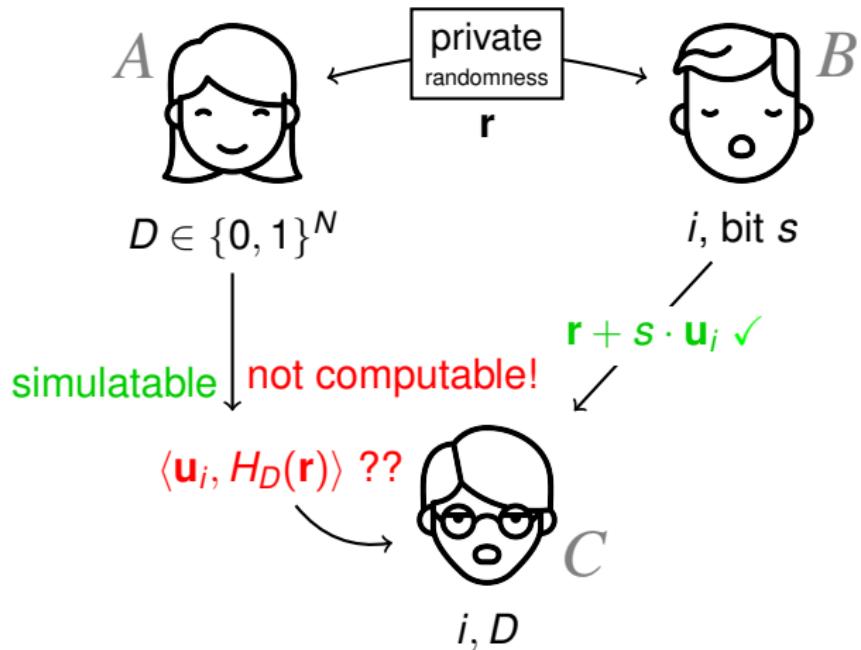
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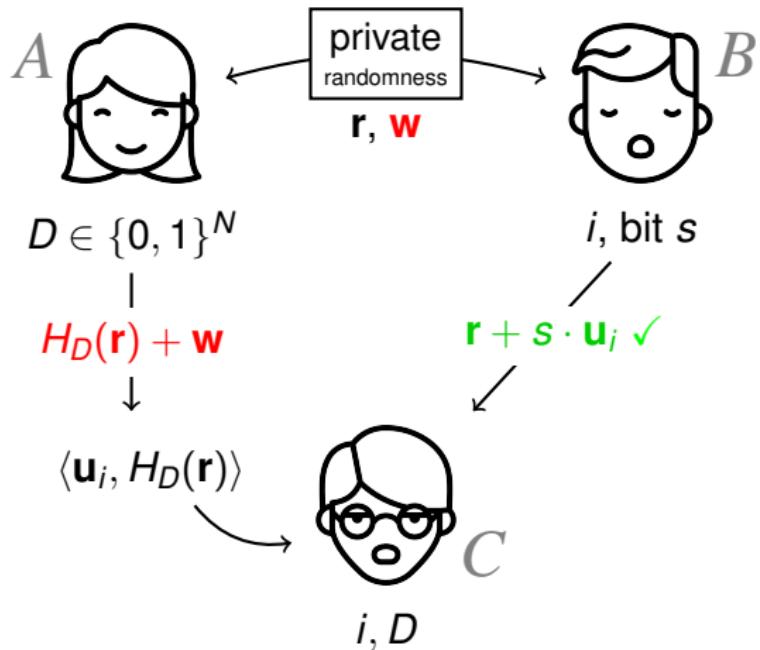
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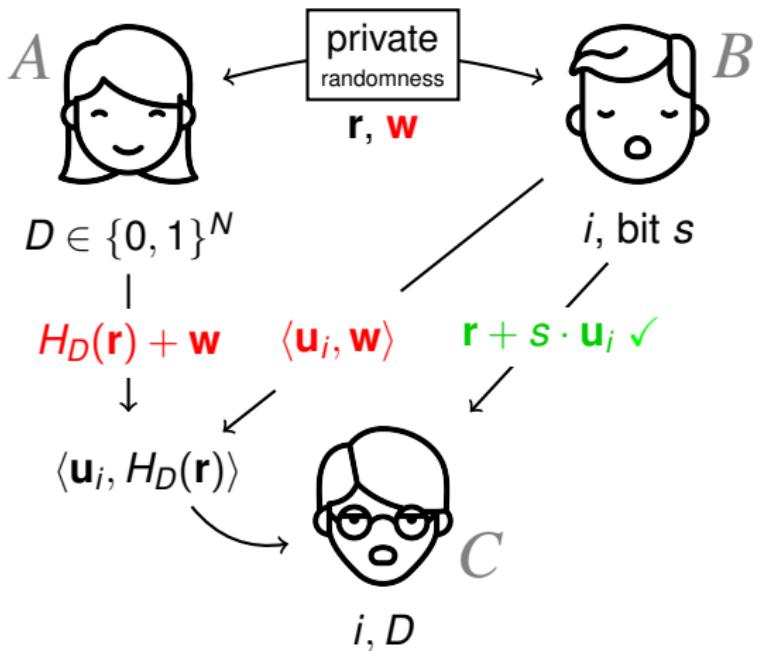
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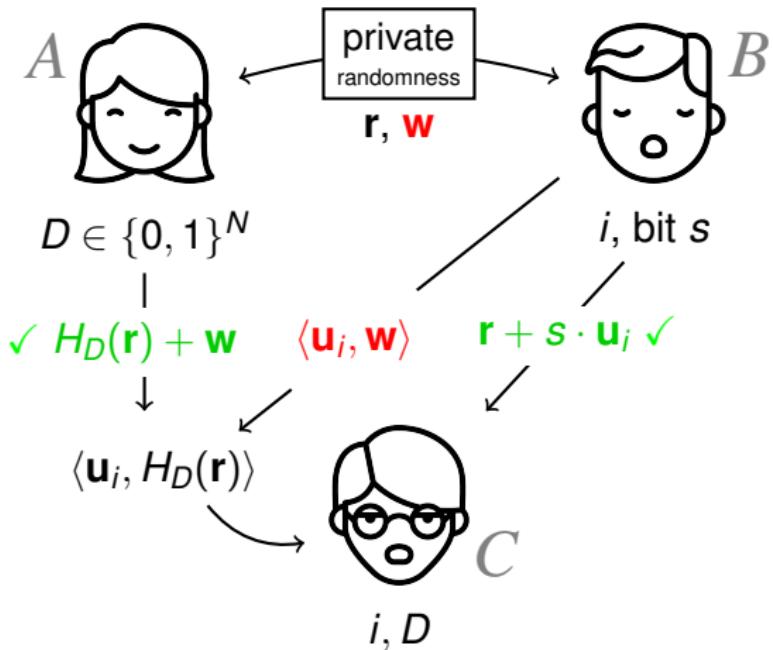
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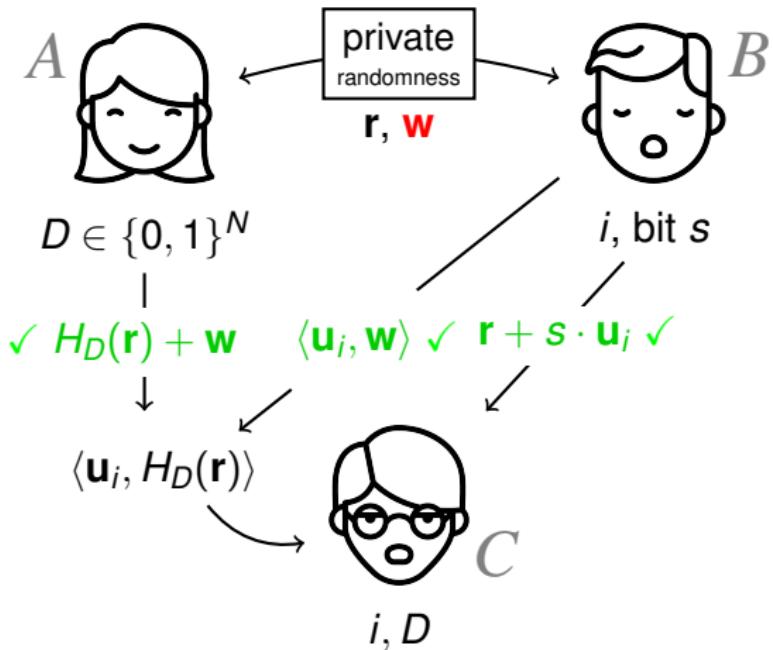
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2-server PIR	CDS
$O(\sqrt{N})$ [CGKS'95] linear server function H_D	
$O(\sqrt[3]{N})$ [CGKS'95, WY'05] quadratic server function H_D	
$2^{\tilde{O}(\sqrt{\log N})}$ [DG'15] general server function H_D	

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2-server PIR	CDS
$O(\sqrt{N})$ [CGKS'95]	$O(\sqrt{N})$
linear server function H_D	linear reconstruction
$O(\sqrt[3]{N})$ [CGKS'95, WY'05]	$O(\sqrt[3]{N})$
quadratic server function H_D	quadratic reconstruction
$2^{\tilde{O}(\sqrt{\log N})}$ [DG'15]	$2^{\tilde{O}(\sqrt{\log N})}$
general server function H_D	general reconstruction

Dvir-Gopi's $\tilde{O}(\sqrt{\log N})$ PIR \implies CDS

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- ▶ Matching vector family $\mathbf{u}_i \in \mathbb{Z}_6^\ell$

$$\langle \mathbf{u}_i, \mathbf{u}_i \rangle = 0 \pmod{6}$$

$$\langle \mathbf{u}_i, \mathbf{u}_j \rangle = 1, 3 \text{ or } 4 \pmod{6} \text{ for } i \neq j$$

and $\ell = 2^{O(\sqrt{\log N \log \log N})}$ [Barrington-Beigel-Rudich'94].

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- ▶ H_D is super non-linear

$$H_D(\mathbf{w}) = \sum_j D_j \cdot \mathbf{u}_j \cdot (-1)^{\langle \mathbf{u}_j, \mathbf{w} \rangle} \pmod{6}$$

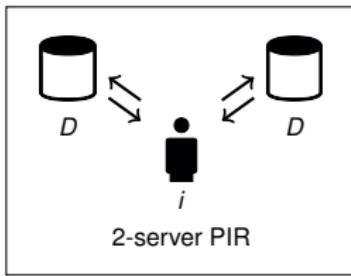
To Summarize

communication complexity	reconstruction
$\Theta(\sqrt{N})$	[GKW'15, ...]
$\Theta(\sqrt[3]{N})$	[This work, GKW'15]
$2^{\tilde{O}(\sqrt{\log N})}$	quadratic
$\Omega(\log N)$	[This work]
	general
	general

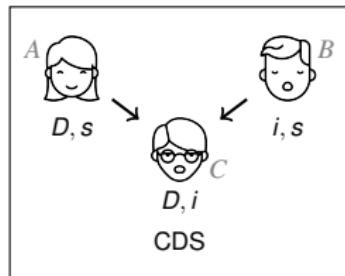
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From



to



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Additional results

Secret sharing on forbidden graph

$O(\sqrt{N})$ [BIKK'14]

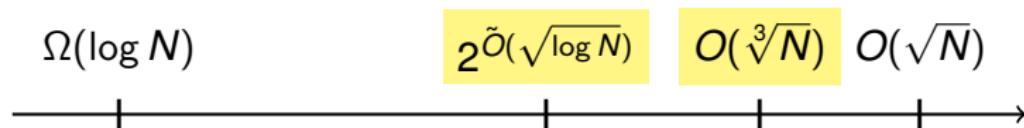
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Open Problems

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$\Theta(\sqrt{N})$	[GKW'15, ...]
$\Theta(\sqrt[3]{N})$	[This work, GKW'15]
$2^{\tilde{O}(\sqrt{\log N})}$	[This work]
$\Omega(\log N)$	[GKW'15]



Open Problems

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$\Theta(\sqrt[3]{N})$	[This work, GKW'15]
$2^{\tilde{O}(\sqrt{\log N})}$	[This work]
$\Omega(\log N)$	[GKW'15]

